

# Social structure and propagation of depositors' panic

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# Motivation

## Seminal paper by Diamond Dybvig JPE (1983)

- Assumptions

- ▶ 1st period depositors invest money to the bank.
- ▶ 2nd period nature reveals their type:
  - ★ Proportion  $s$  of them are impatient and withdraw money immediately.
  - ★ Proportion  $1 - s$  are patient and play coordination game.

- Conclusions

- ▶ Deposits can provide allocation superior to those of exchange market.
- ▶ There is a multiple equilibria, one of which is always a bank run.
- ▶ Government provision of insurance may produce superior outcome

# Motivation

- Depositors decisions are partially sequential
  - ▶ Descriptions of bank runs Sprague (1910), Wicker (2001)
  - ▶ Statistical data Starr and Yilmaz (2007)
- Many depositors make decision observing actions of others:
  - ▶ Kelly and Grada (2000) - bank run of Turkey's Islamic financial houses in 2001.
  - ▶ Iyer and Puri (2008) consider depositor level data for a bank that faced a run in India in 2001.
- Main contribution: we introduce social network as a coordination mechanism that depositors may use to make their decision.

# Model

- There is a continuum of agents.
- Agents are embedded into the network of personal contacts, represented by a random graph with degree distribution  $p(k)$ .
- At period 0 each agent invests 1 unit into a bank account.
- At period 1 nature reveals agents type in 2 steps:
  - ▶ Nature draws proportion of impatient agents  $s$  in the society from distribution with CDF  $Q(s)$ .
  - ▶ According to realized  $s$  nature assigns types to depositors.

# Model

- Impatient depositor withdraws money regardless of prevailing conditions.
- Patient depositor with  $k$  links withdraws according to the strategy  $P_w(m, k) \in [0, 1]$ .
- If a depositor withdraws money from the bank she gets pay-off  $a(w(s))$
- If depositor waits till 2nd period she gets pay-off  $b(w(s))$
- We assume single crossing property of  $b(w) - a(w)$ :  $b(w) > a(w)$  for  $w < \bar{w}$  and  $b(w) < a(w)$  for  $w > \bar{w}$

# Model

- Probability that a randomly chosen neighbor of depositor withdraws:

$$\hat{w} = s + (1 - s) \sum_{k=1}^{\infty} \xi(k) \sum_{m=0}^{k-1} P_w(m, k) \frac{(k-1)!}{m!(k-1-m)!} \hat{w}^m (1 - \hat{w})^{k-1-m},$$

where  $\xi(k)$  is the degree distribution of depositor's neighbor.

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- Proportion of agents that withdraw is:

$$w(s, \hat{w}) = s + (1 - s) \sum_{k=0}^{\infty} p(k) \sum_{m=0}^k P_w(m, k) \frac{k!}{m!(k-m)!} \hat{w}^m (1 - \hat{w})^{k-m}$$

# Maximization problem

- Depositor solves the following maximization problem:

$$\sum_{m=0}^k P(M = m|k) \int_0^1 [(1 - P_w(m, k))b(w(s)) + P_w(m, k)a(w(s))] P(S = s|m, k) ds$$

- $P(M = m|k)$  is the probability to observe  $m$  out of  $k$  neighbors withdrawing
- $P(S = s|m, k)$  Bayesian updating of belief about true state  $s$ .



# Optimal decision

## Proposition

*Optimal decision strategy of agent  $P_w(m, k)$  is a cut-off rule, such that agent withdraws if  $m \geq m_k$  and waits otherwise. Moreover, cut-off value  $m_{k+1} \in \{m_k, m_{k+1}\}$ .*

# Maximization problem

- Knowing that optimal strategy is cut-off rule maximization problem becomes:

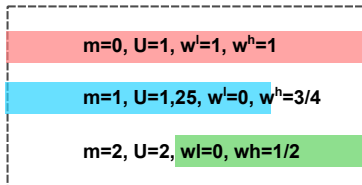
$$\int_0^{\bar{w}} b(w) dQ(w^{-1}(s)) + \int_{\bar{w}}^1 a(w) dQ(w^{-1}(s)) - \\ - \int_0^{\bar{w}} [b(w) - a(w)] l_w(m_k, k+1 - m_k) dQ(w^{-1}(s)) - \\ - \int_{\bar{w}}^1 [a(w) - b(w)] [1 - l_w(m_k, k+1 - m_k)] dQ(w^{-1}(s))$$

- The second term is loss due to the 1st type error (false positive)
- The third term is 2nd type error (false negative)

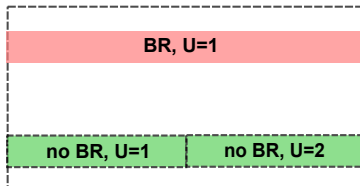
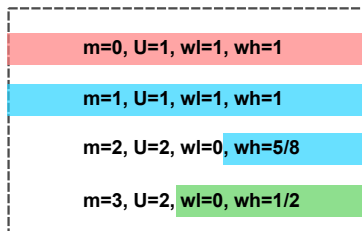
# Examples

- States:  $s_l = 0, s_h = \frac{1}{2}, q = \frac{1}{2}$  Pay-off:  $a(w) = 1, b(w)$  equals 2 for  $w < \bar{w}$  and 0 otherwise.

**k=1**



**k=2**



# Optimal decision

## Proposition

*Depositor's utility is increasing function in the number of links.*

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*Depositor's utility is increasing function in the number of links.*

- Assume that it is optimal for the agent to have  $\frac{m_k}{k} = \frac{1}{2}$ .
- Depositor with 2 links by setting  $m_k = 1$  has exactly the optimal cut-off value.
- Depositor with 3 links can approximate optimal cut-off strategy only by  $\frac{1}{3}$  or  $\frac{2}{3}$ .

# Continuous Case

## Proposition

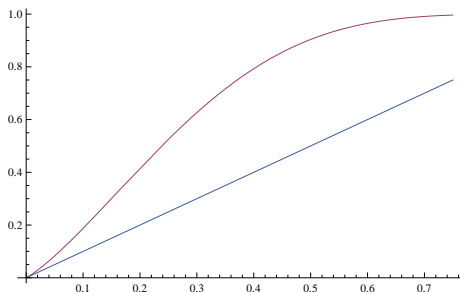
*For an arbitrary degree distribution and exogenously given cut-off rule  $m_k = \alpha k$ , if mean degree converges to infinity then the following holds:*

$$\hat{w} = \begin{cases} s, & s < \alpha \\ 1, & \text{otherwise} \end{cases}$$

# Continuous Case

## Proposition

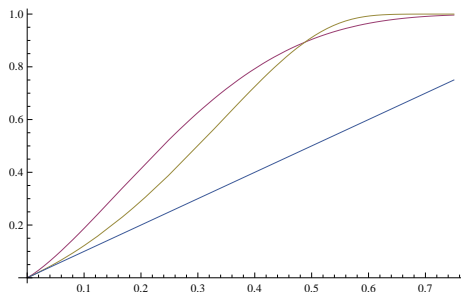
For any two degree distributions  $F(k)$  and  $\tilde{F}(k)$  and for their corresponding neighboring node degree distributions  $G(k)$  and  $\tilde{G}(k)$ , if the following holds:  $\tilde{F}(k)$  **FOSD**  $F(k)$  and  $\tilde{G}(k)$  **FOSD**  $G(k)$  and for any  $k$ ,  $0 < m(k+1) - m(k) < 1$ , then there are  $\underline{s}$  and  $\bar{s}$ , such that  $w_{\tilde{F}}^*(s) < w_F^*(s)$  for  $s < \underline{s}$  and  $w_{\tilde{F}}^*(s) > w_F^*(s)$  for  $s > \bar{s}$ .



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